Wishes for the VeriPB proof format

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- SAT, MAXSAT, PBO solvers
- Main development between 2004 and 2011
- Specific work by Emmanuel Lonca (Multi-Objective Optimization in 2015) and Romain Wallon (PB proof systems in 2020)
- Contains PB solvers with either Resolution-based or Cutting-Planes-based proof systems

- 2013: DRUP UNSAT proof (Daniel)
- 2021: VeriPB 1 UNSAT proof (Antony Blomme and Romain Wallon)
- 2024: iDRUP incremental proof and VeriPB 2 optimal and UNSAT (Daniel)

- UNSAT: one call to the solver (can have restarts)
- Incremental: multiple calls to the SAT solver, log calls under assumptions.
- Optimal proof: multiple calls to the SAT solver

Sat4j solvers emit events, which can be used for proof logging

DRUP proof logging: easy

public void end(Lbool result) {
 if (result == Lbool.FALSE) {
 out.println("0"); out.close(); }}

public void learn(lConstr c) { printConstr(c);}

```
public void delete(IConstr c) {
        out.print("d "); printConstr(c); }
```

```
private void printConstr(IConstr c) {
  for (var i = 0; i < c.size(); i++) {
    out.print(LiteralsUtils.toDimacs(c.get(i)));
    out.print(" ");
} out.println("0");}</pre>
```

```
public void learnUnit(int p) {
    out.print(p); out.println(" 0");}
```

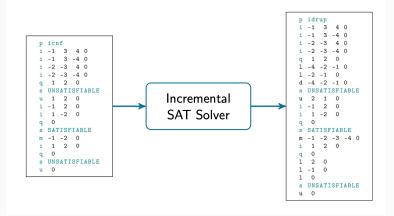
```
public void withReason(PBConstr constr) {
    if (constr != null) {
        this.reason = new StringBuilder("" + constr.getId());}}
public void weakenOnReason(int p) {
    this.reason.append(" x" + Math.abs(p)).append(" w");}
public void weakenOnReason(BigInteger coeff, int p) {
reason.append(" " + coeff).append(" x" + Math.abs(p))
            .append(" W');}
public void weakenOnConflict(int p) {
    this.conflict.append(" x" + Math.abs(p)).append(" w");}
```

```
public void weakenOnConflict(BigInteger coeff, int p) {
    this.conflict.append(" " + coeff).append(" x" + Math.abs(p))
        .append(" W");}
```

- Proposed in Katalin Fazekas, Florian Pollitt, Mathias Fleury, Armin Biere: Certifying Incremental SAT Solving. LPAR 2024: 321-340
- Implemented in Sat4j during the master thesis of Iris Parruca studying API fuzzing in Sat4j, advised by Mathias Fleury and Tobias Paxian
- Log all events of incremental SAT: learning clauses, queries with assumptions, models, unsat assumptions.
- Log several search spaces, either SAT or UNSAT
- icnf format and idrup-checker available

https://github.com/arminbiere/idrup-check

icnf and idrup example [Fazekas et al 2024]



https://kfazekas.github.io/talks/mbmv2024_talk.pdf

- Some solvers in Sat4j can derive clauses but not "watch" them: MiniLearning in 2004
- Those clauses may be used in conflict analysis, but do not propagate nor cause conflicts
- Those clauses needed to be logged in the proof
- New "ignore" event created

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That problem would have been hard to notice without API fuzzing

- Benefit from iDRUP implementation for handling multiple calls to the SAT solver
- Passes the test cases from VeriPB 2.0 repository
- VeriPB 1.0 to 2.0 changes were under-estimated



Olivier Roussel [PB24] output DERIVABLE IMPLICIT À : Daniel Le Berre

Salut Daniel,

Pourquoi utilises-tu dans tes preuves "output DERIVABLE IMPLICIT" et pas "output NONE" ?

A+ OR



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Certificates were accepted by VeriPB but not by CakePB

VeriPB 2 proof logging: the output in Sat4j

```
private void displayConclusion (String answer) throws IOException {
    fw.write("output DERIVABLE IMPLICIT\n"):
   fw.write("conclusion " + answer + "\n");
public void endWhole(ExitCode status) {
        switch (status) {
        case UNSATISFIABLE:
            if (!this.foundContradiction) {
                fw.write("rup \geq 1 \ n");
                this.nConstraints++:
            if (((IPBSolverService) solverService)
                    .getObjectiveFunction() != null) {
                displayConclusion("BOUNDS INF INF");
            } else {
                displayConclusion ("UNSAT");
            break:
        case OPTIMUM_FOUND:
            displayConclusion ("BOUNDS " + lowerBound + " " + upperBound);
            break:
        case SATISFIABLE:
            displayConclusion ("SAT"):
            break:
        default:
            fw.write("output NONE\n");
            fw.write("conclusion NONE\n");
        fw.write("end pseudo-Boolean proof"):
        fw.close():
```

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- PB competition benchmarks contain "unit clauses"
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Sat4j Resolution VeriPB 2.0 certificates are correct on those benchmarks

- "Unit clauses" are propagated directly when parsing the benchmark
- The simplification is performed to represent the constraint in the most appropriate way in the solver, again while parsing the benchmark
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VeriPB 2.0 support does not ship with Sat4j, it lives in a specific branch (VERIPB2)

On Sat4j side:

- create new events for all simplifications occurring before the search?
- not so easy on 20 years old code (47k LOC)
- API fuzz testing can help (not available for PB yet in Sat4j)

On VeriPB side:

- Could VeriPB be more friendly with equivalent transformations?
- Could VeriPB focus on what we derive, not how we derive it?